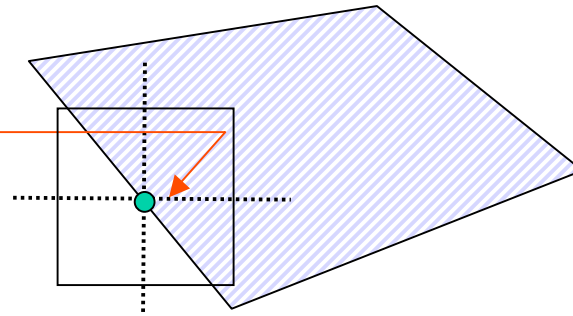


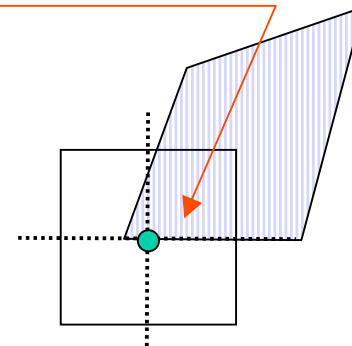
# Ambiguous cases

- What if a pixel is exactly on the edge?
- Polygons are usually adjacent to other polygons, so we want a **convention** which will give the pixel to *one* of the adjacent polygons or the *other* (as much as possible).
- “Draw left and bottom edges”

- if  $(x+\delta, y)$  is in, pixel is in  
(for sufficiently small, positive,  $\delta$ )



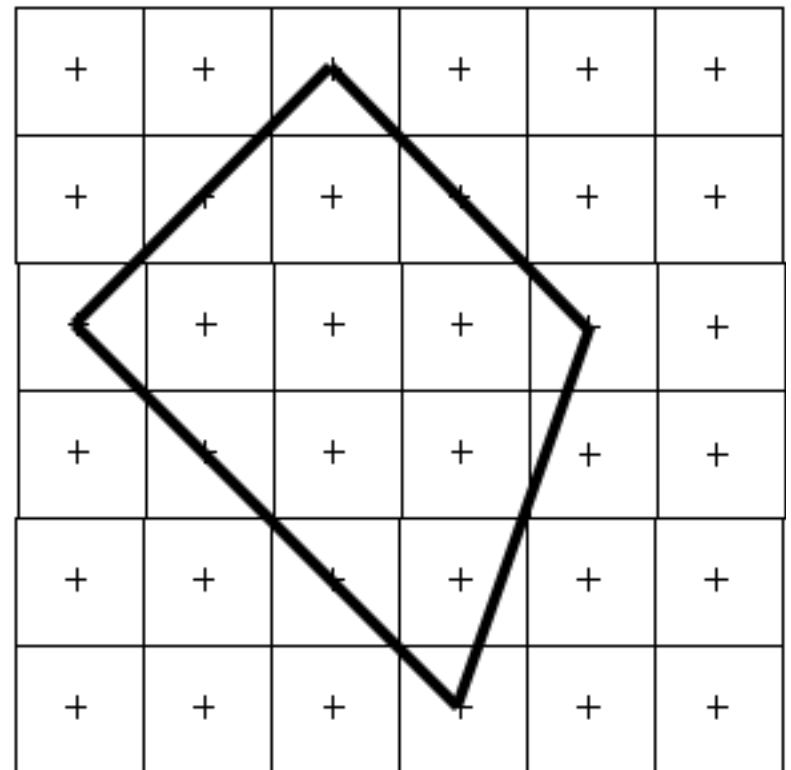
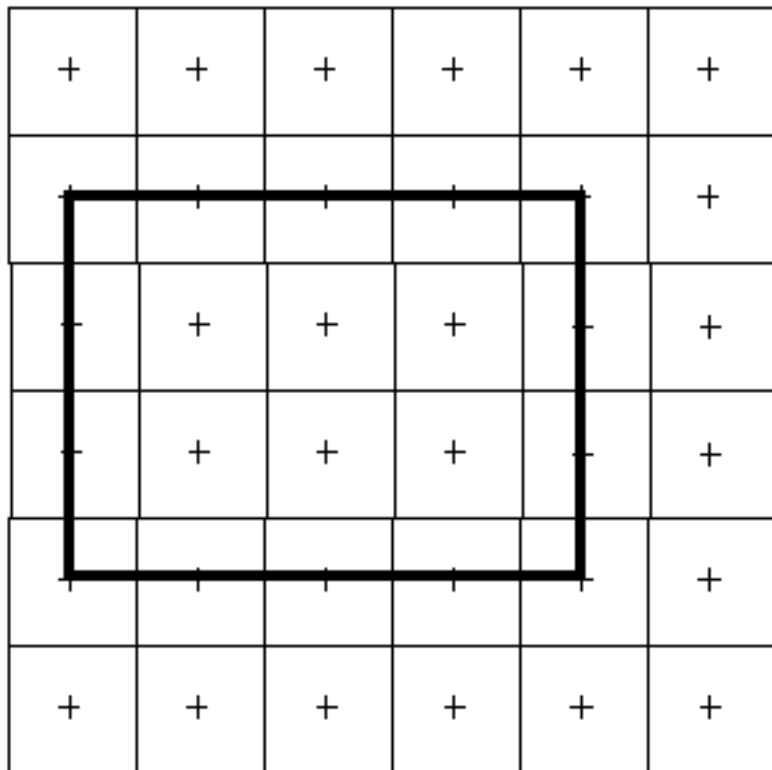
- horizontal edge? if  $(x+\delta, y+\epsilon)$  is in, pixel is in  
(for sufficiently small, positive,  $\delta, \epsilon$ )



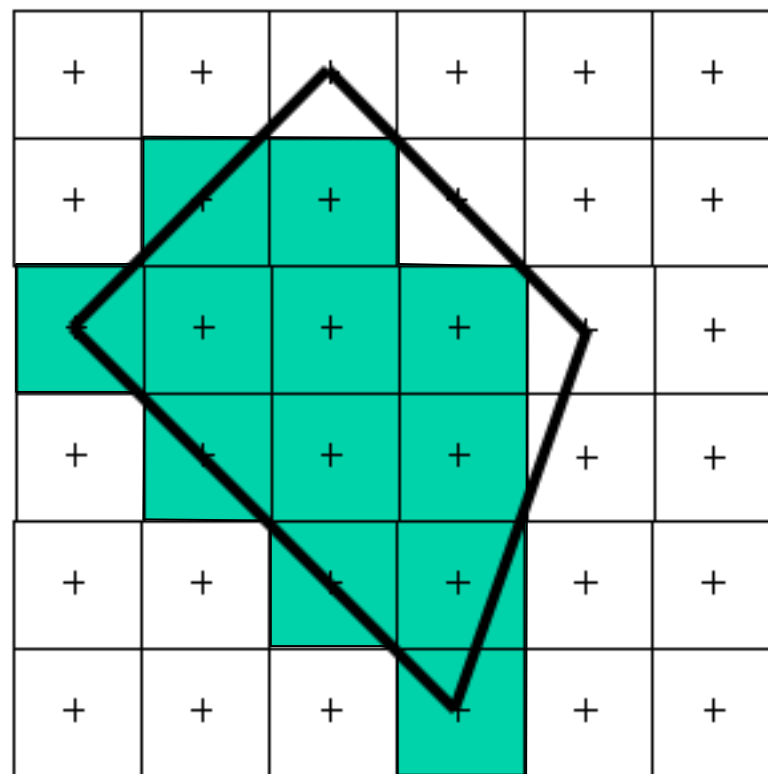
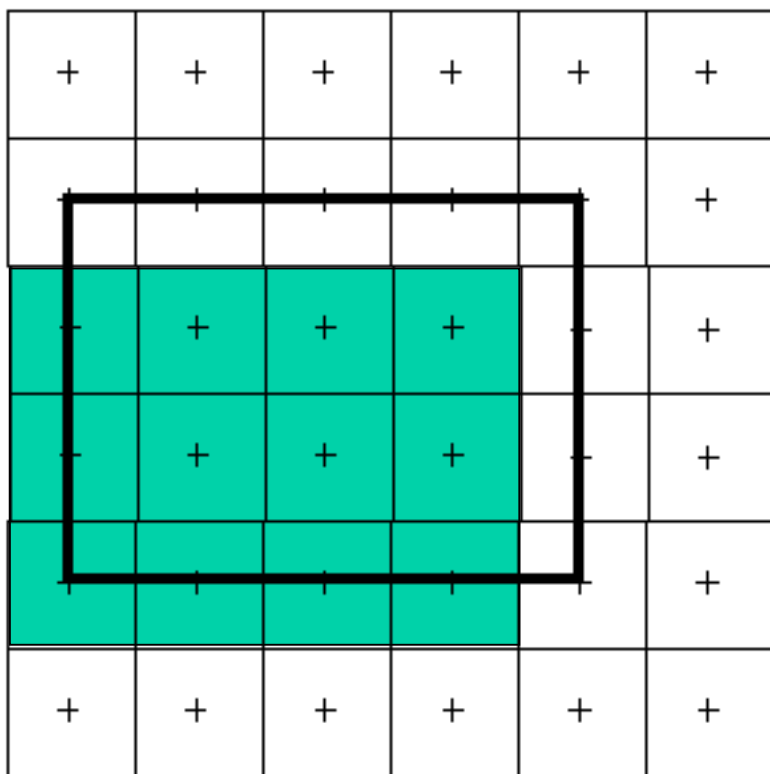
## Ambiguous cases (2)

- Vertex?--essentially draw those on “left and bottom”, but detailed analysis should be done in conjunction with one’s scan conversion algorithm

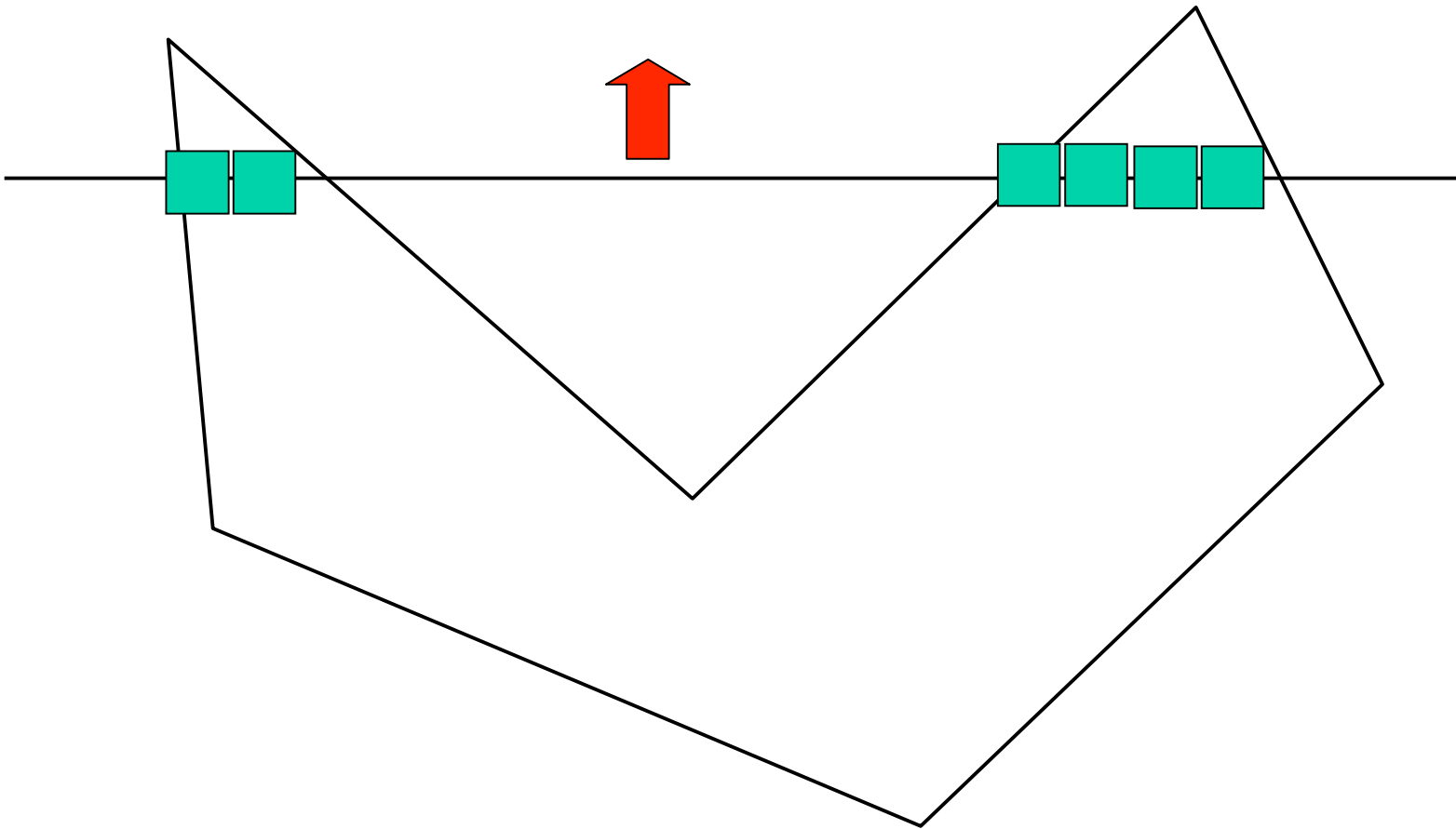
# Ambiguous inside cases (?)



## Ambiguous inside cases (answer)



# Sweep fill



# Sweep fill

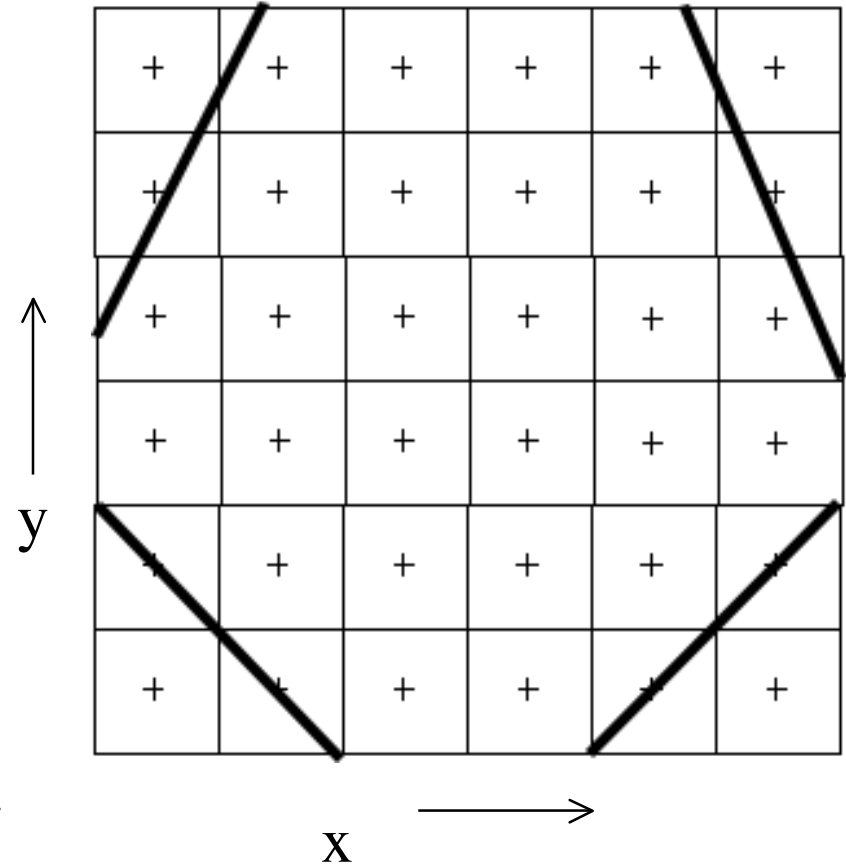
- Reduces to filling many spans
- Inside/outside parity is relatively straightforward
- Need to compute the spans, then fill
- Need to update the spans for each scan
- Need to implement “inside” rule for ambiguous cases.

# Spans

- Process - fill the bottom horizontal span of pixels; move up and keep filling
- Suppose we have the span, i.e., have  $x_{min}$ ,  $x_{max}$  designating the intersection of the span with the sweep line (floating point).
- Need pixels with integers greater than **or** equal to  $x_{min}$ , but less than  $x_{max}$ . (Less than **or** equal  $x_{max}$  would break the convention)
- In code

for ( $x = \text{ceil}(x_{min})$ ;  $x < \text{ceil}(x_{max})$ ;  $x++$ )

recall that  $\text{ceil}(x)$  is largest integer greater than or equal to  $x$



- Note that two adjacent polygons do not fight each for any pixels

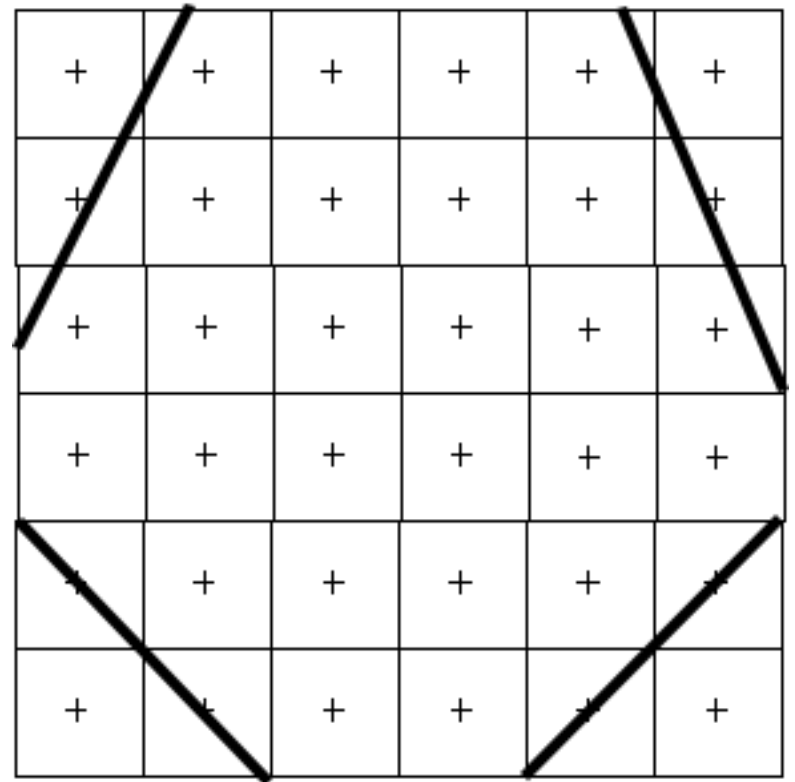
# Algorithm

- For each row in the polygon:
  - Throw away newly irrelevant edges
  - Obtain newly relevant edges
  - Fill spans
  - Update spans
- Issues:
  - What aspects of edges need to be stored?
  - When is an edge relevant/irrelevant?



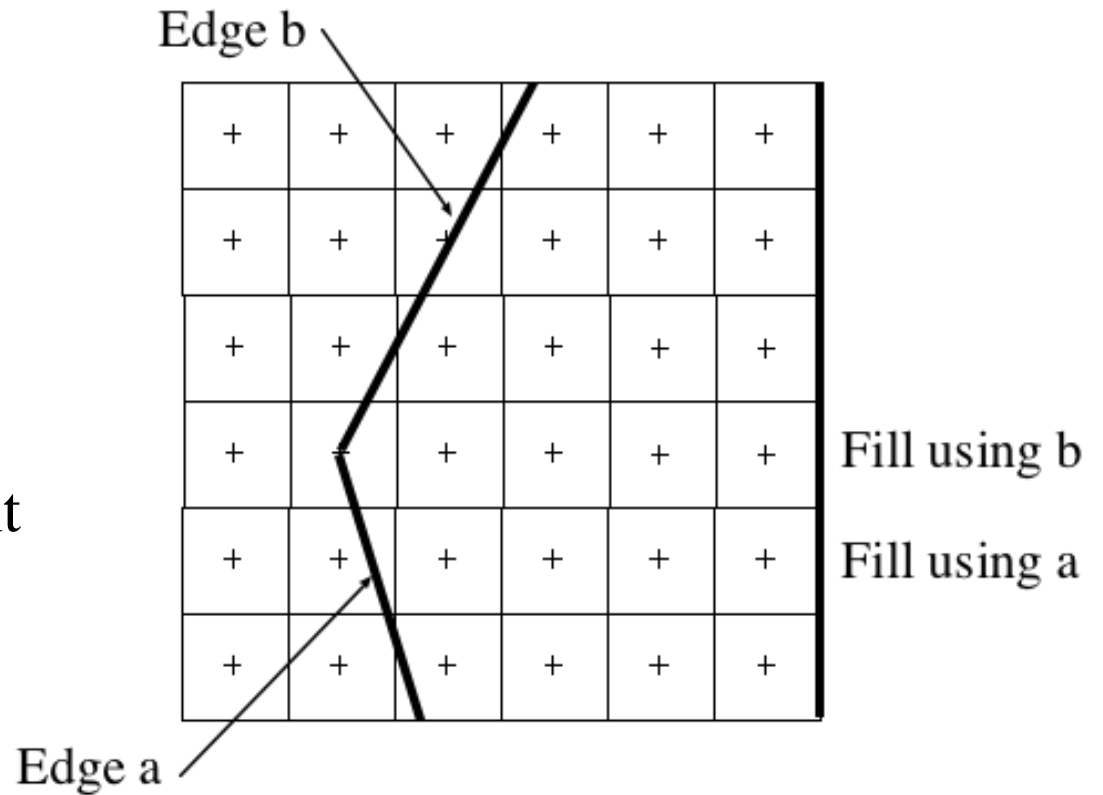
# The next span - 1

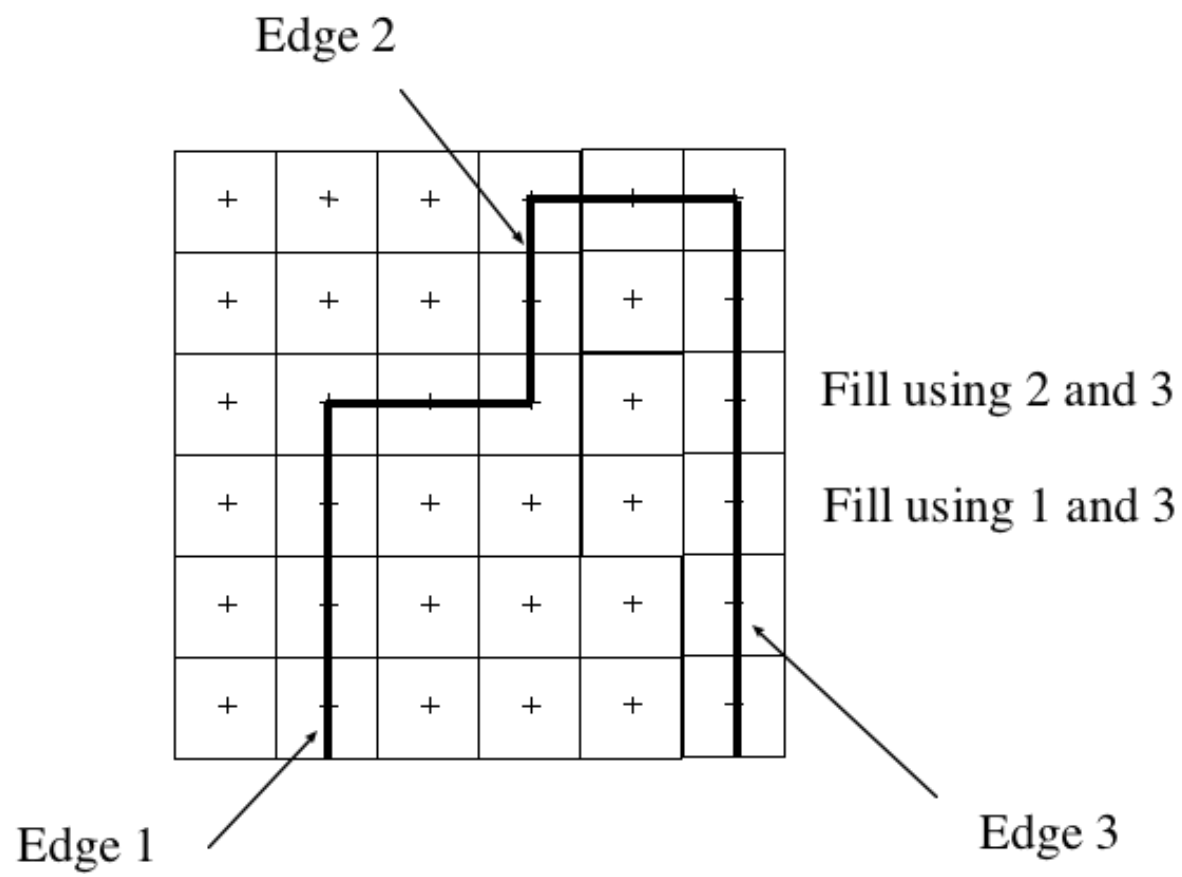
- For an edge, have  $y=mx+b$
- hence, if  $y_n=m x_n +b$ , then  $y_{n+1}=y_n+1=m (x_n+1/m)+b$
- Hence, *if there is no change in the edges*, we have:  
     $x_{\max} \rightarrow x_{\max}+(1/m)(x_{\max})$   
     $x_{\min} \rightarrow x_{\min}+(1/m)(x_{\min})$



## The next span - 2

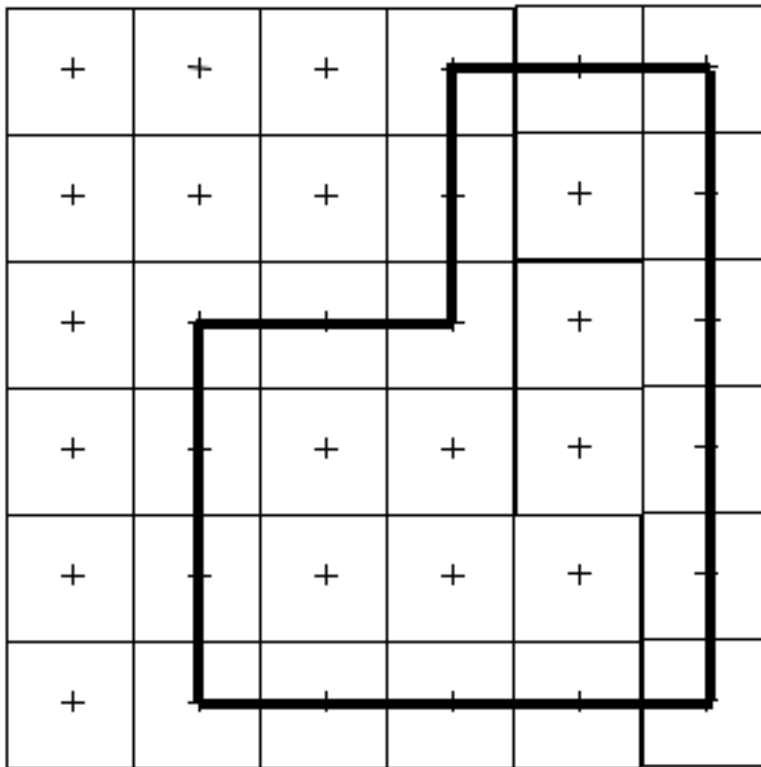
- Horizontal edges are irrelevant
- Edge is irrelevant when  $y \geq y_{\max}$  of edge (note appeal to convention)
- Similarly, edge is relevant when  $y \geq y_{\min}$  of edge





# Filling in details

- Maintain a list of active edges in case there are multiple spans of pixels (Active Edge List).
- For each edge on the list, must know: x-value (current, initialize to min), max y value of edge,  $1/m$
- Keep edges in a table, indexed by minimum y value - Edge Table
- For row = min to row=max
  - AEL=append(AEL, ET(row));
  - remove edges whose  $y_{max}=row$
  - sort AEL by x-value
  - fill spans
  - update each edge in AEL



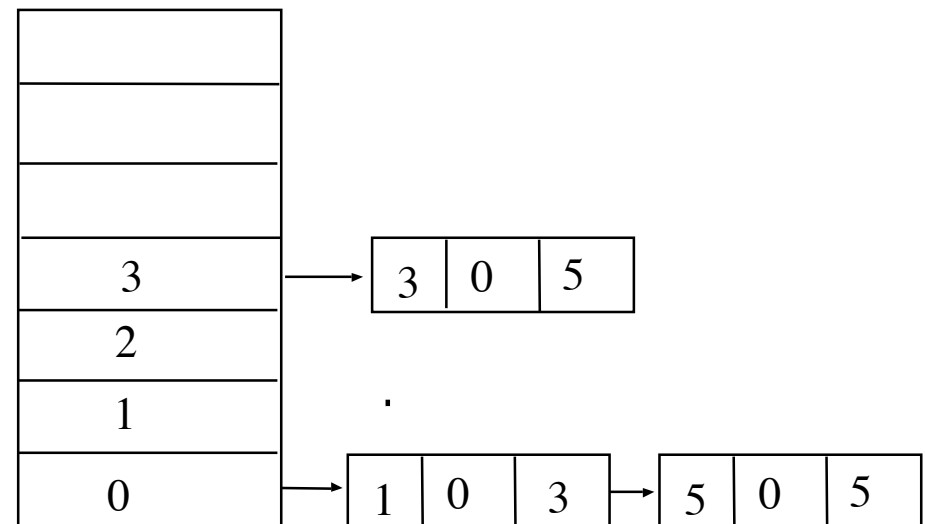
## Example one

Compute the edge table (ET\_  
to begin. Then fill polygon and  
update active edge list (AEL)  
row by row. (Ignore horizontal edges)

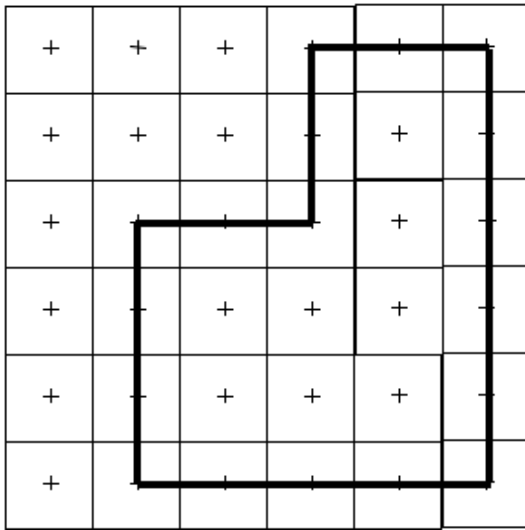
Format of AEL entries

xmin	1/m	ymin
------	-----	------

ET



The AEL entries begin  
with xmin, and are  
initialized at row ymin

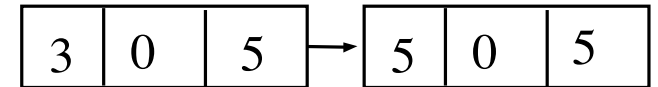


## Example one

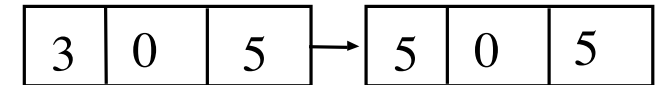
AEL just before filling listed row

Row=5

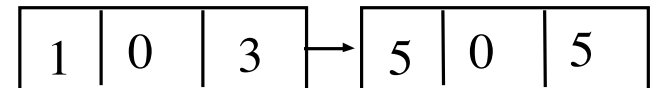
Row=4



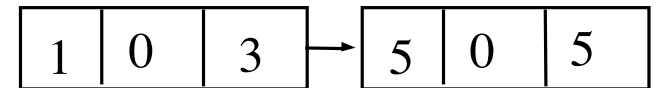
Row=3



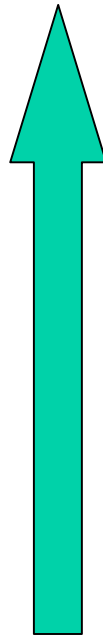
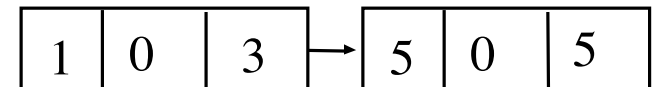
Row=2



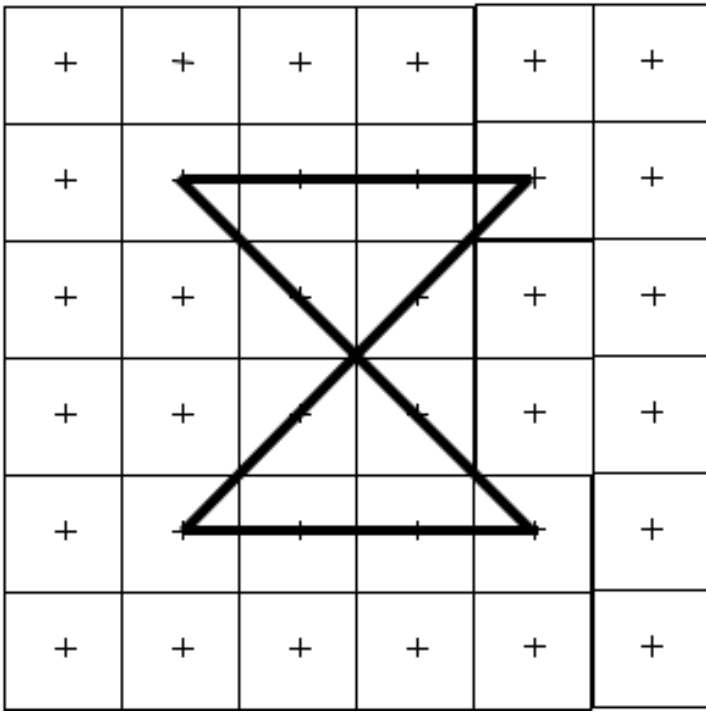
Row=1



Row=0



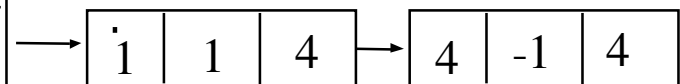
## Example two



The AEL entries begin with xmin, and are initialized at row ymin

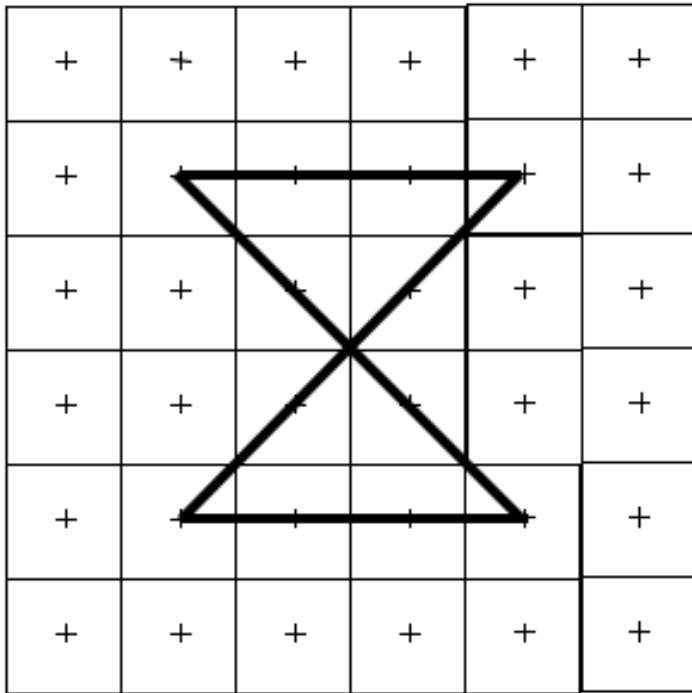
ET

4
3
2
1
0



Format of AEL entries

xmin	1/m	ymin
------	-----	------

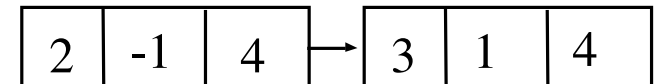


## Example two

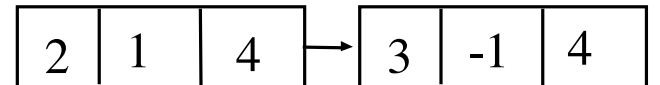
AEL just before filling  
listed row

Row=4

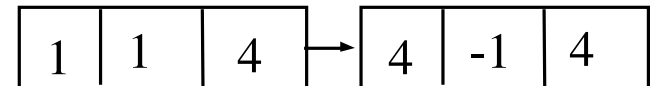
Row=3



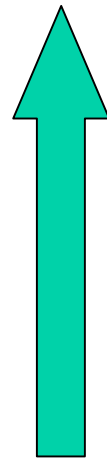
Row=2



Row=1



Row=0





# Comments

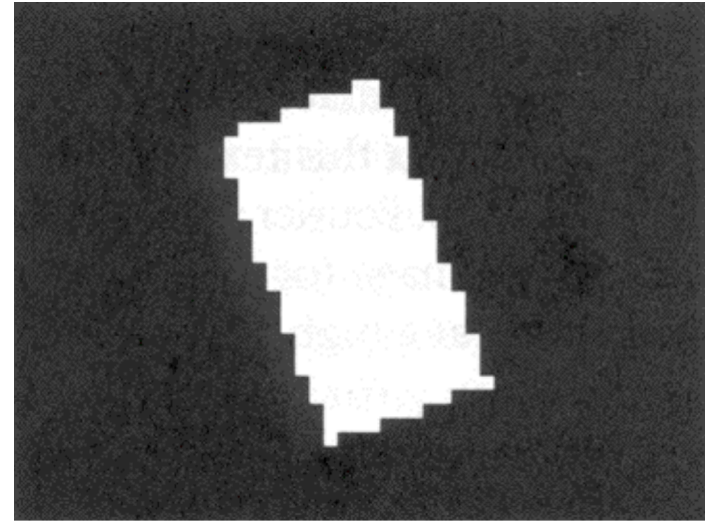
- Sort is quite fast, because AEL is usually almost in order.
- Nonetheless, OpenGL limits to convex polygons, so two and only two elements in AEL at any time, and no sorting.
- With additional logic to keep track of what color to use, can fill in many polygons at a time.
- Can be done *without* floating point

# Dodging floating point

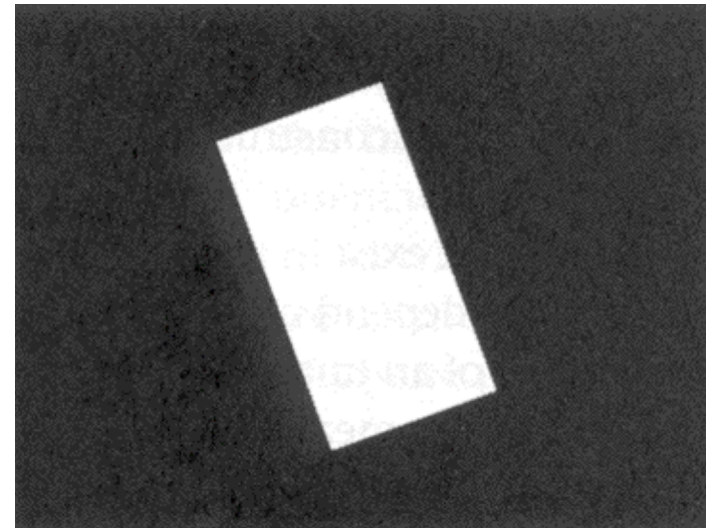
- For edge,  $1/m = Dx/Dy$ , which is a rational number.
- Store  $x$  as  $x\_int$ ,  $x\_num$ ,  $x\_denom = Dy$
- then  $x \rightarrow x + 1/m$  is given by:
  - $x\_num = x\_num + Dx$
  - if  $x\_num \geq x\_denom$ 
    - $x\_int = x\_int + 1$
    - $x\_num = x\_num - x\_denom$
- Advantages:
  - no floating point
  - can tell if  $x$  is an integer or not (check  $x\_num = 0$ ), and get  $\text{truncate}(x)$  easily, for the span endpoints.

# Aliasing/Anti-Aliasing

- Analogous to the case of lines
- Anti-aliasing is done using graduated gray levels computed by by smoothing and sampling
- Problem with “slivers” is really an aliasing problem.



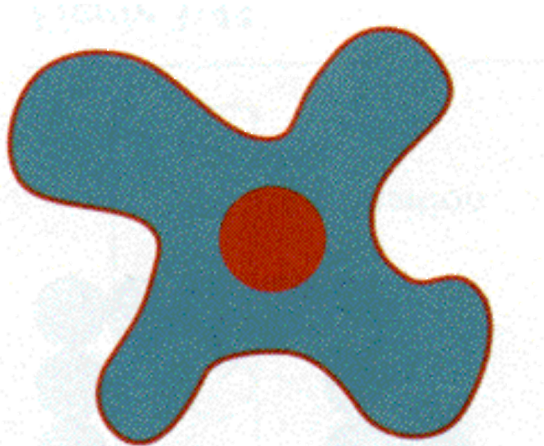
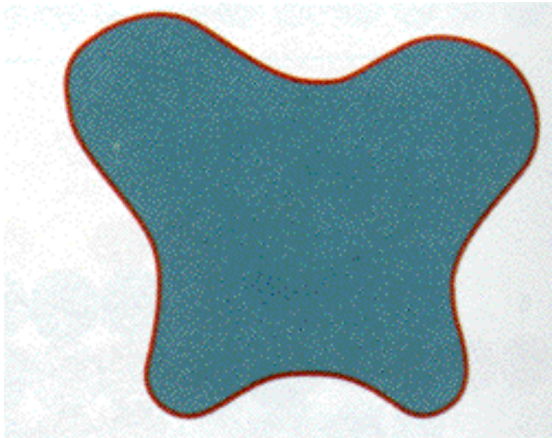
Aliasing



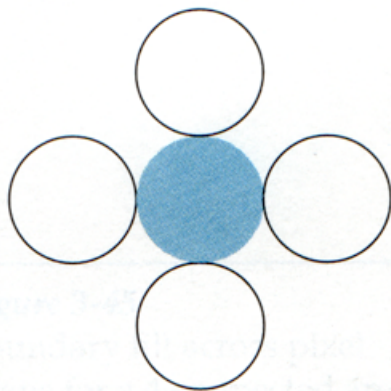
Ideal

# Boundary fill

- Basic idea: fill in pixels inside a boundary
- Recursive formulation:
  - to fill starting from an inside point
  - if point has not been filled,
    - fill
    - call on all neighbours that are not boundary pixels.



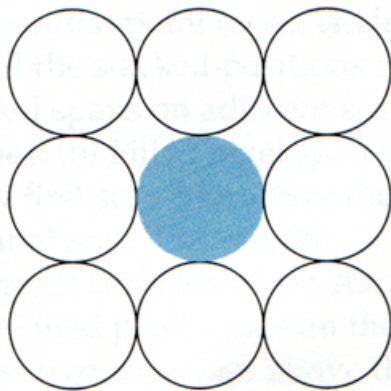
# Choice of neighbours is important



(a)

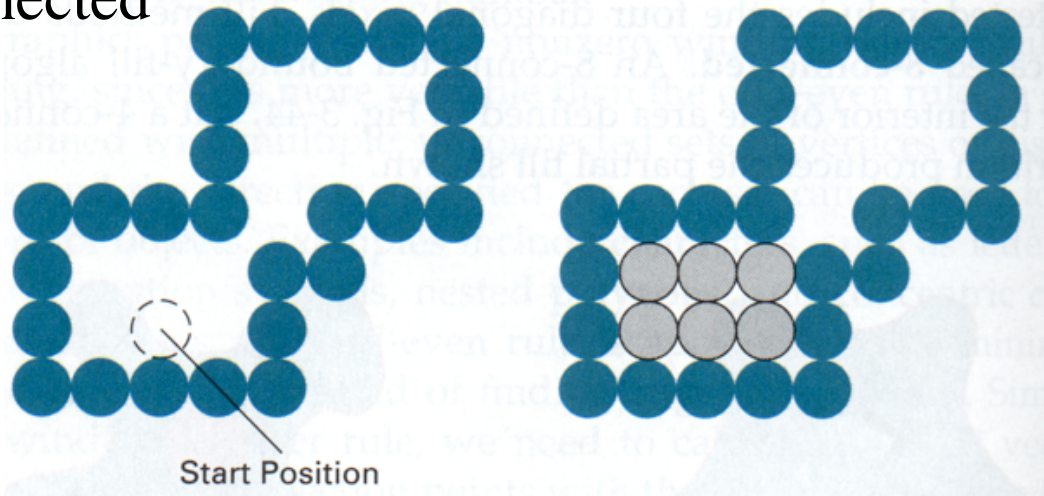
4-connected

4 connected fill of  
a four connected  
boundary doesn't work



(b)

8 connected



Start Position

- Filled Pixel Spans**

**Stacked Positions**

2
1

3
1

6
5
4
1

5
4
1

# Pattern fill

- Use index into screen as index into pattern