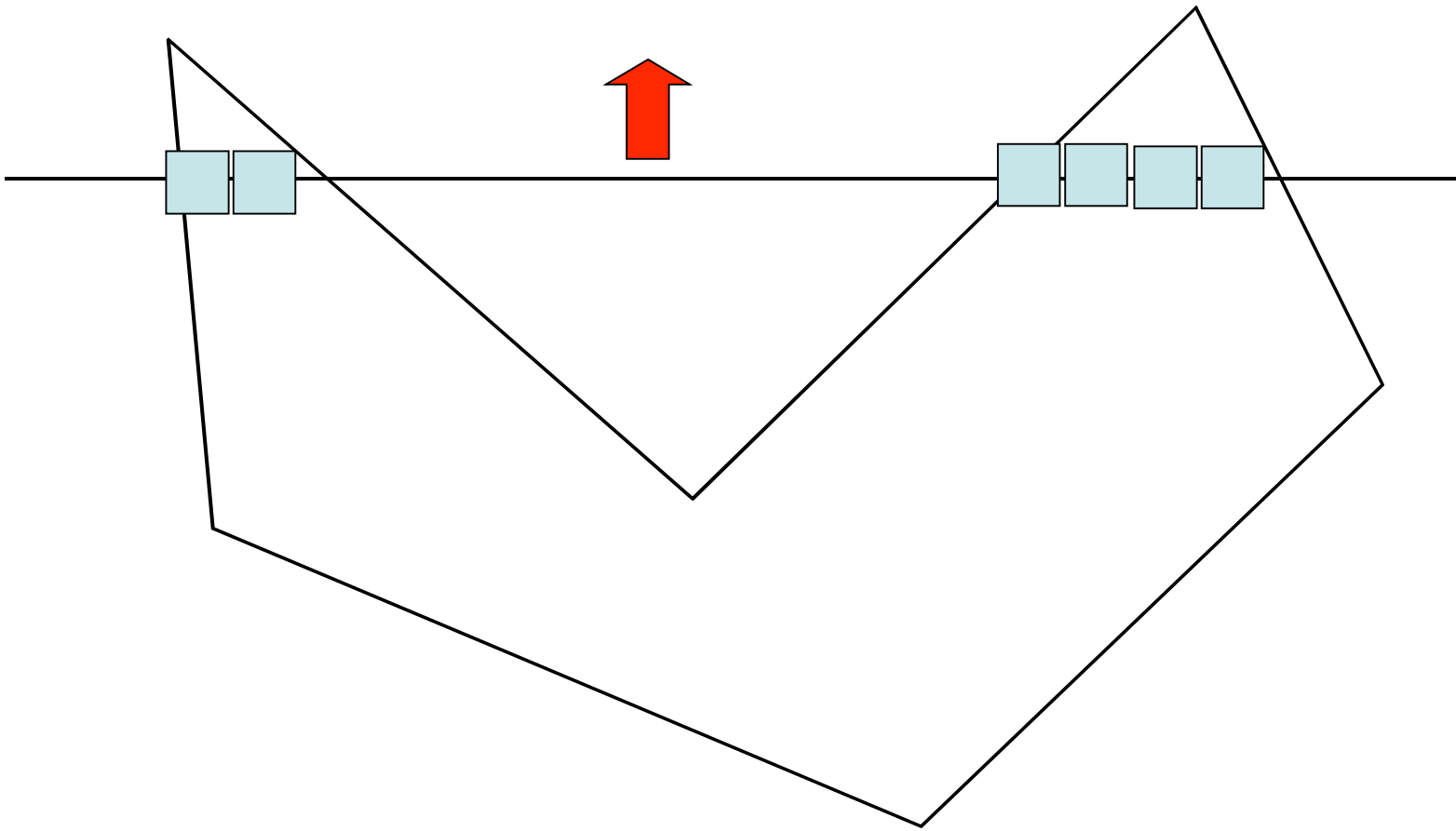


Sweep fill

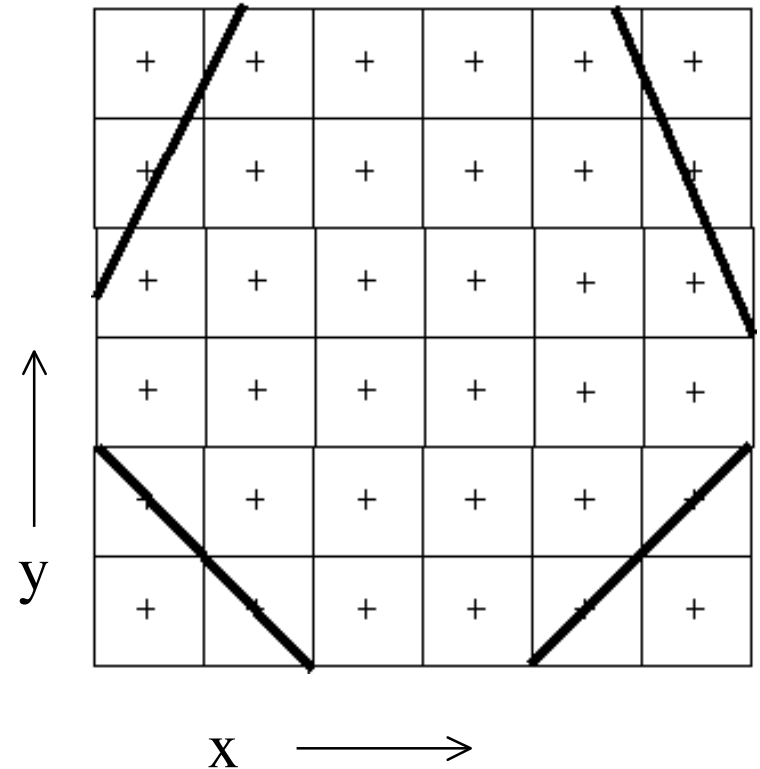


Sweep fill

- Reduces to filling many spans
- Inside/outside parity is relatively straightforward
- Need to compute the spans, then fill
- Need to update the spans for each scan
- Need to implement “inside” rule for ambiguous cases.

Spans

- Fill the bottom horizontal span of pixels; move up and keep filling
- Assume we have x_{min} , x_{max} .
- Recall--for non integral x_{min} (going from outside to inside), **round up** to get first interior point, for non integral x_{max} (going from inside to outside), **round down** to get last interior point
- Recall--convention for for integral gives a span closed on the left and open on the right
- **Thus:** fill from $\text{ceiling}(x_{min})$ up to but not including $\text{ceiling}(x_{max})$

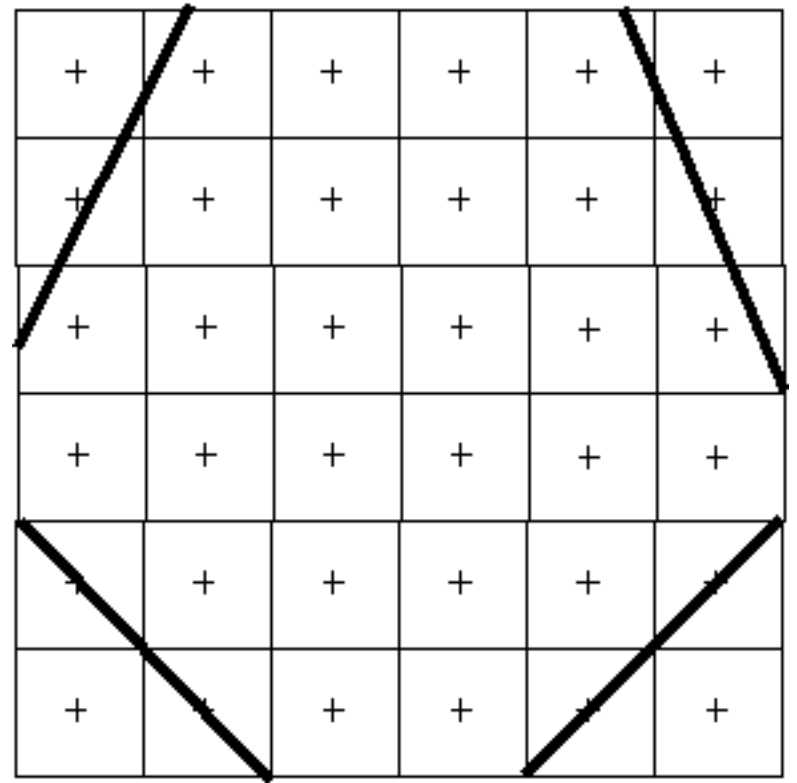


Algorithm

- For each row in the polygon:
 - Throw away irrelevant edges (horizontal ones, ones that we are done with)
 - Obtain newly relevant edges (ones that are starting)
 - Fill spans
 - Update spans
- Issues:
 - what aspects of edges need to be stored?
 - when is an edge relevant/irrelevant?

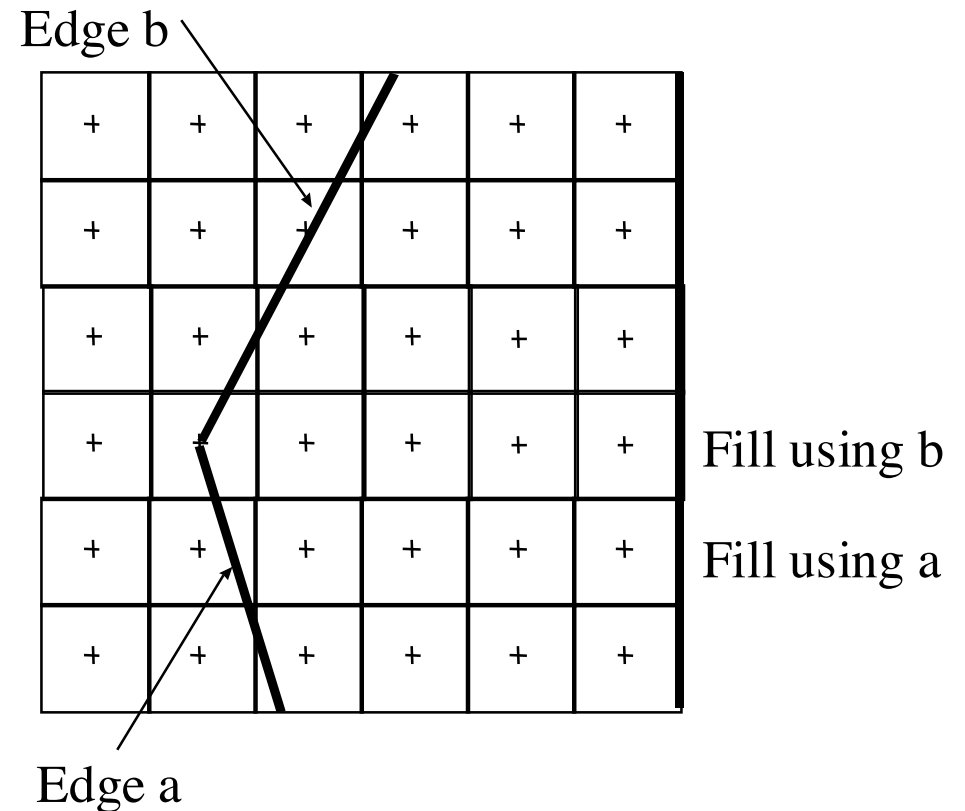
The next span - 1

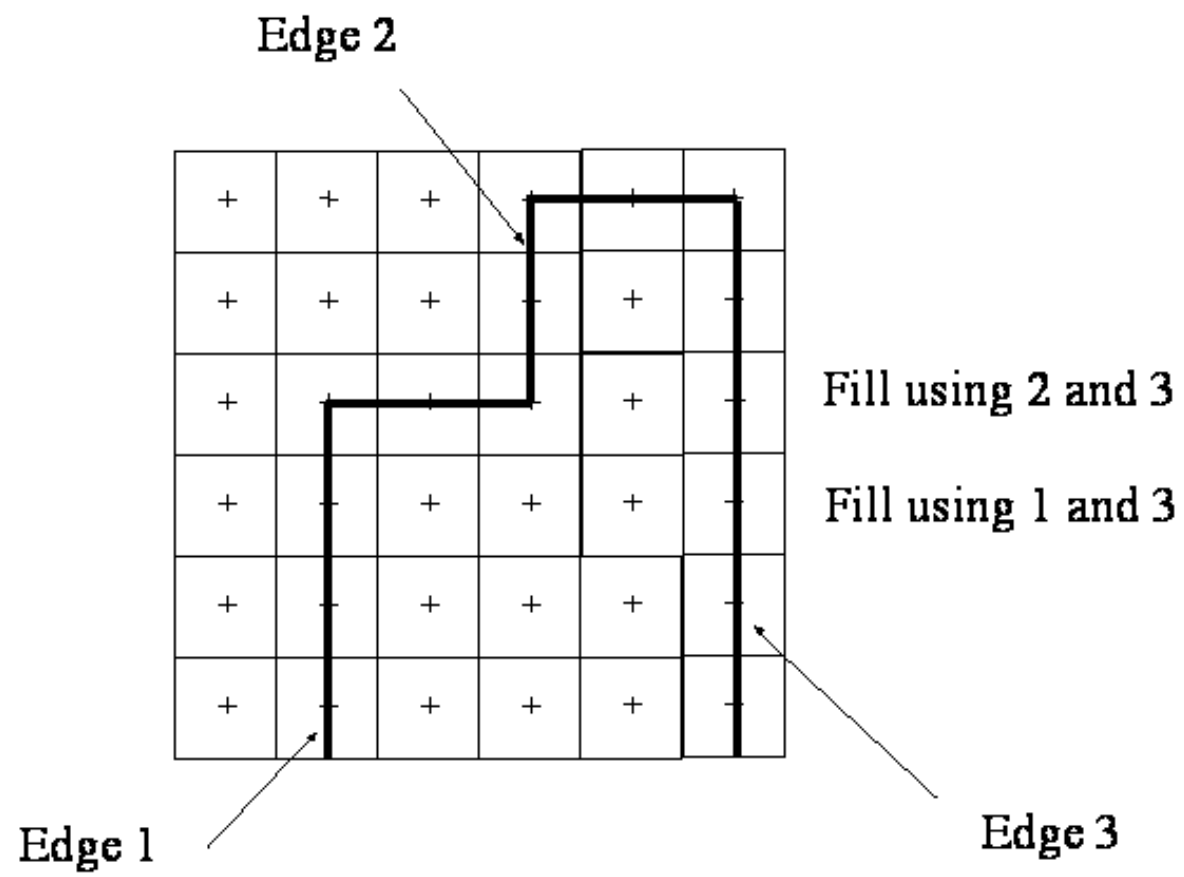
- for an edge, have $y=mx+c$
- hence, if $y_n=m x_n +c$, then $y_{n+1}=y_n+1=m (x_n+1/m)+c$
- hence, *if there is no change in the edges*, have:
 $x += (1/m)$



The next span - 2

- Horizontal edges are irrelevant (typically would be pruned at the outset)
- Edge becomes relevant when $y \geq y_{\min}$ of edge
- Edge becomes irrelevant - when $y \geq y_{\max}$ of edge (note appeal to convention)



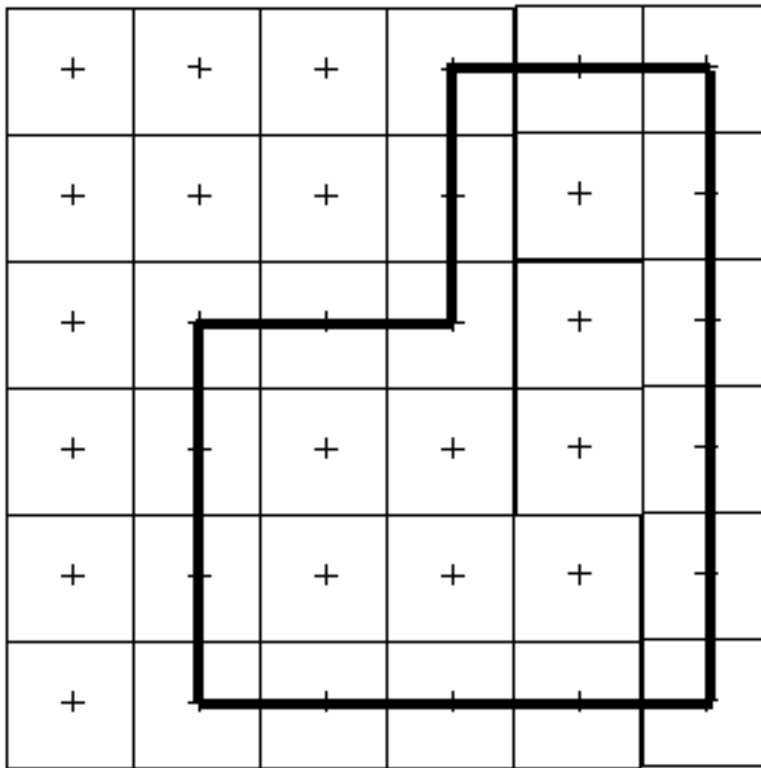


Filling in details -- 1

- For each edge store: x-value, maximum y value of edge, $1/m$
 - x-value starts out as x value for y_{\min}
 - m is never 0 because we ignore horizontal ones
- Keep edges in a table, indexed by minimum y value (Edge Table==ET)
- Maintain a list of active edges (Active Edge List==AEL).

Filling in details -- 2

- For row = min to row=max
 - AEL=append(AEL, ET(row)); (add edges starting at the current row)
 - remove edges whose ymax=row
 - OK since we are assuming integral coordinates; otherwise one would use ceil(ymax)
 - sort AEL by x-value
 - fill spans
 - parity rule
 - convention for integral x_{\min} and x_{\max}
 - update each edge in AEL
 - $x += (1/m)$

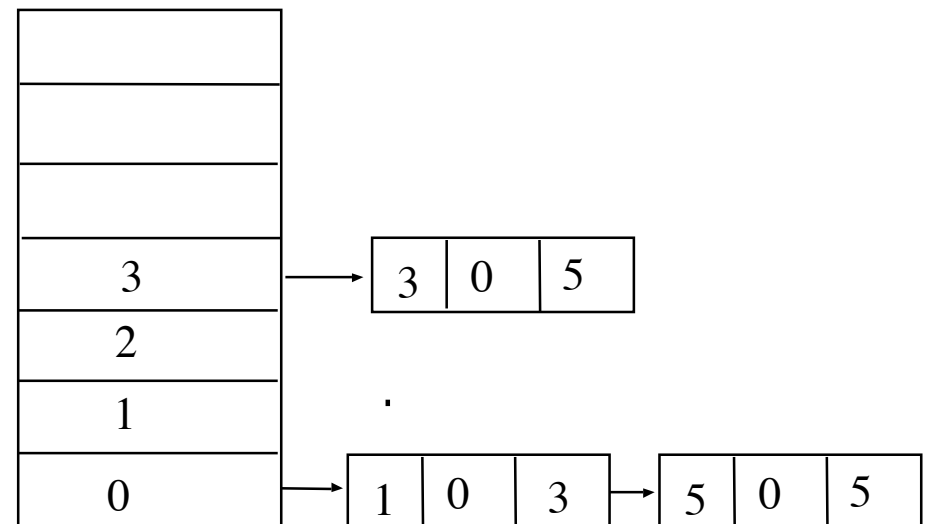


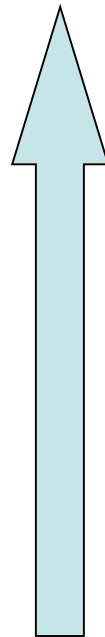
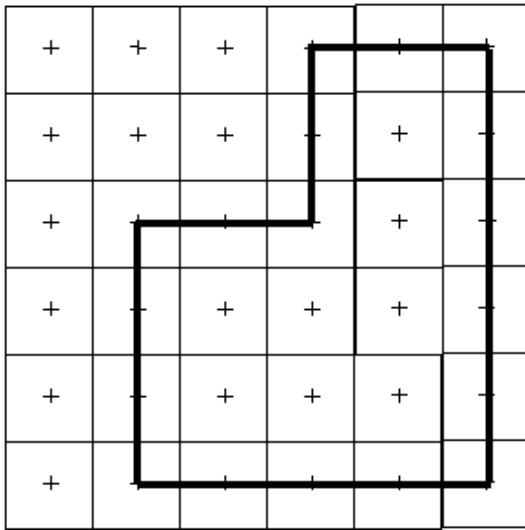
Compute the edge table (ET) to begin. Then fill polygon and update active edge list (AEL) row by row.

Format of edge entries

| x | 1/m | y _{max} |
|---|-----|------------------|
|---|-----|------------------|

ET





AEL just before filling

Row=5

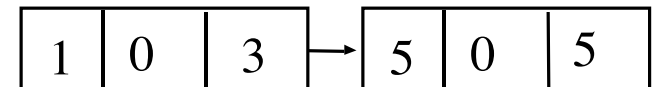
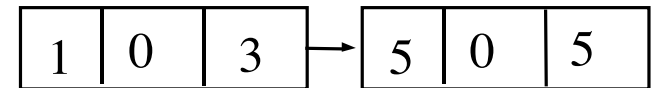
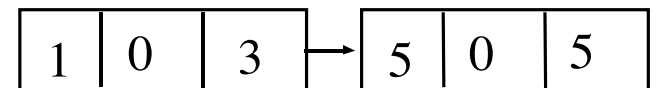
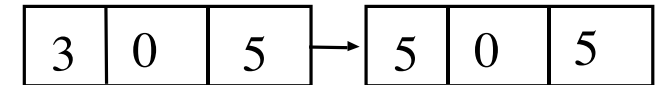
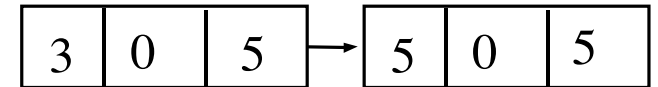
Row=4

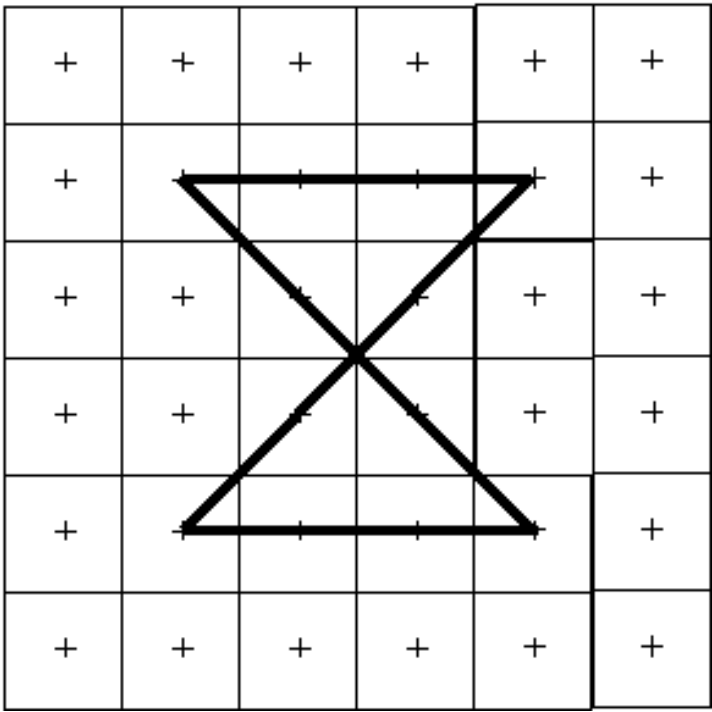
Row=3

Row=2

Row=1

Row=0



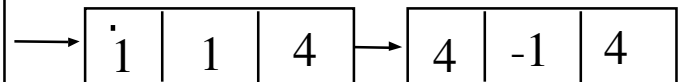


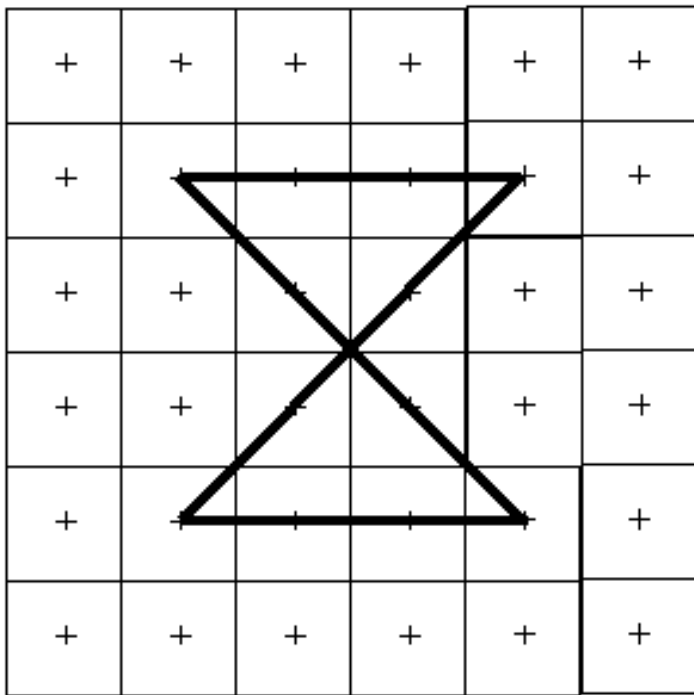
Format of edge entries

| x | 1/m | y _{max} |
|---|-----|------------------|
|---|-----|------------------|

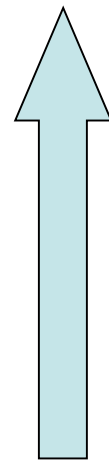
ET

| |
|---|
| |
| |
| 4 |
| 3 |
| 2 |
| 1 |
| 0 |



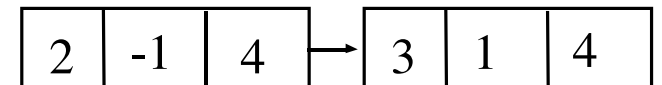


AEL just before filling

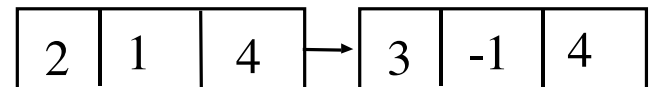


Row=4

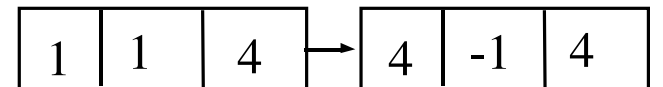
Row=3



Row=2



Row=1



Row=0

Comments

- Sort is quite fast, because AEL is usually almost in order.
- Nonetheless, OpenGL limits to convex polygons, so two and only two elements in AEL at any time, and no sorting.
- With additional logic to keep track of what color to use, can fill in many polygons at a time.
- Can be done *without* floating point

Dodging floating point

- $1/m = Dx/Dy$, which is a rational number.
- $x = x_int + x_num/Dy$
- store x as (x_int, x_num) ,
- then $x \rightarrow x + 1/m$ is given by:
 - $x_num = x_num + Dx$
 - if $x_num \geq x_denom$
 - $x_int = x_int + 1$
 - $x_num = x_num - x_denom$
- Advantages:
 - no floating point
 - can tell if x is an integer or not (check $x_num=0$), and get $\text{truncate}(x)$ easily, for the span endpoints.